

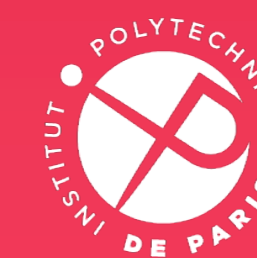
Data Protection and Privacy in a Quantum World

Sébastien Canard

Journées Nationales du GDR

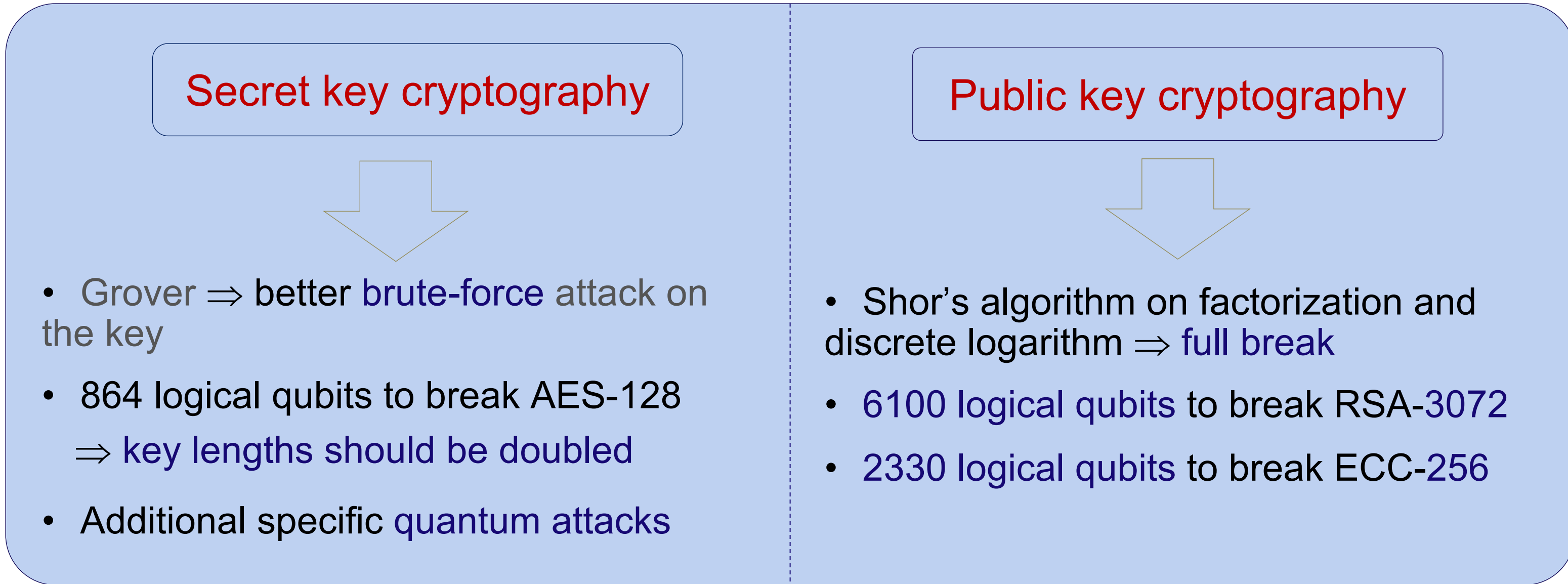
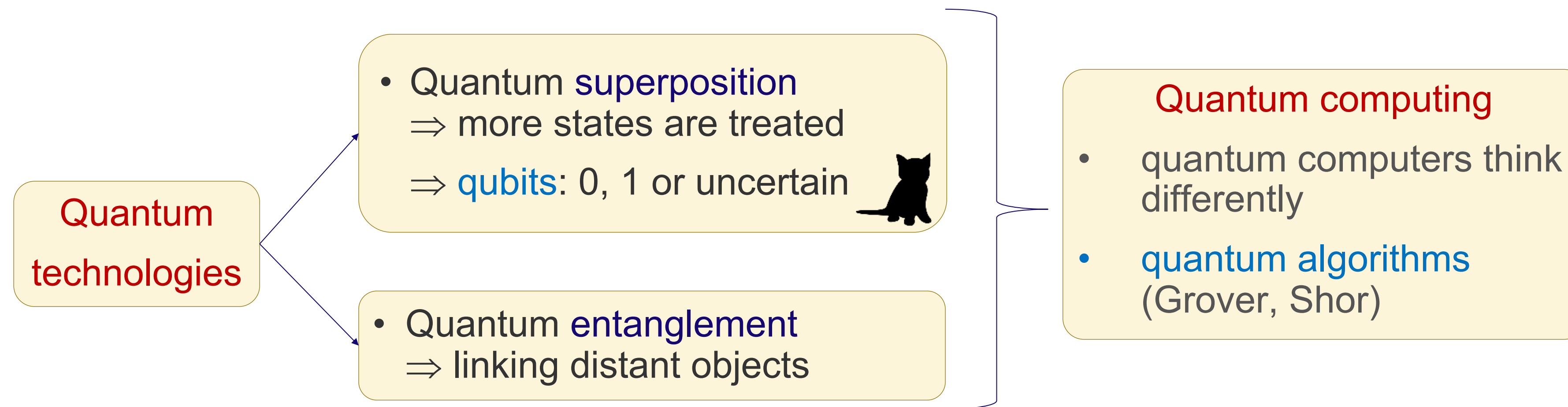
Sécurité Informatique

June 2023

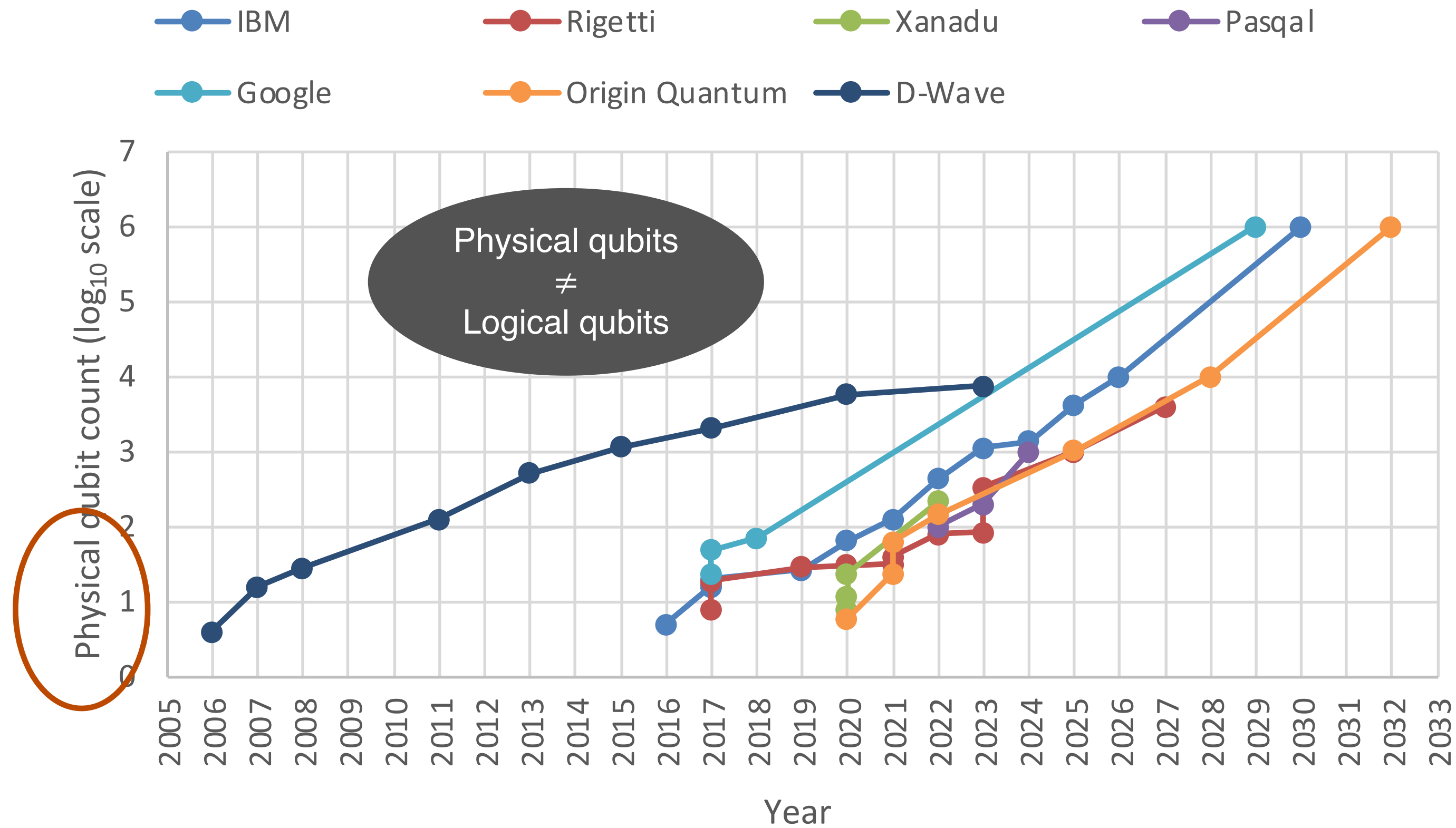


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Quantum computers and cryptography



Availability of quantum computers



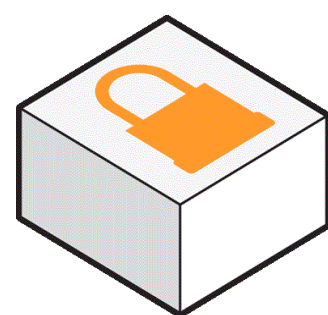
Not a threat now!

But « store now, decrypt later » paradigm

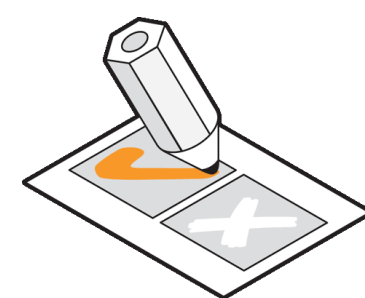
Recall:

AES	864 logical qubits
RSA	6100 logical qubits
ECC	2330 logical qubits

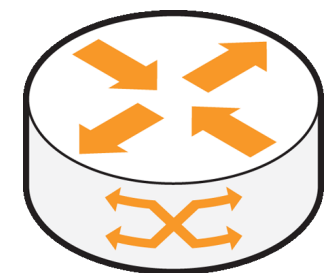
Impacted primitives



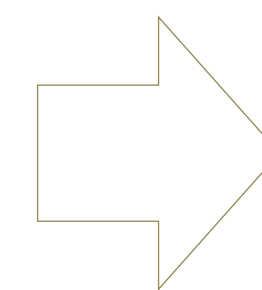
- **KEM and Encryption**
 - Data confidentiality
 - Using additional secret key cryptography or not



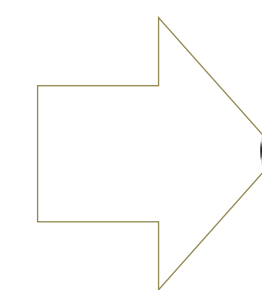
- **Digital signatures**
 - Person/message authentication
 - Integrity and non-repudiation



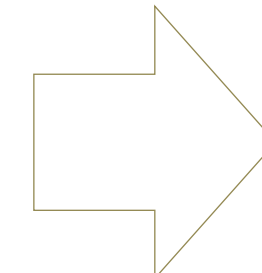
- **Advanced cryptography**
 - Privacy-preserving techniques
 - Sensitive data protection



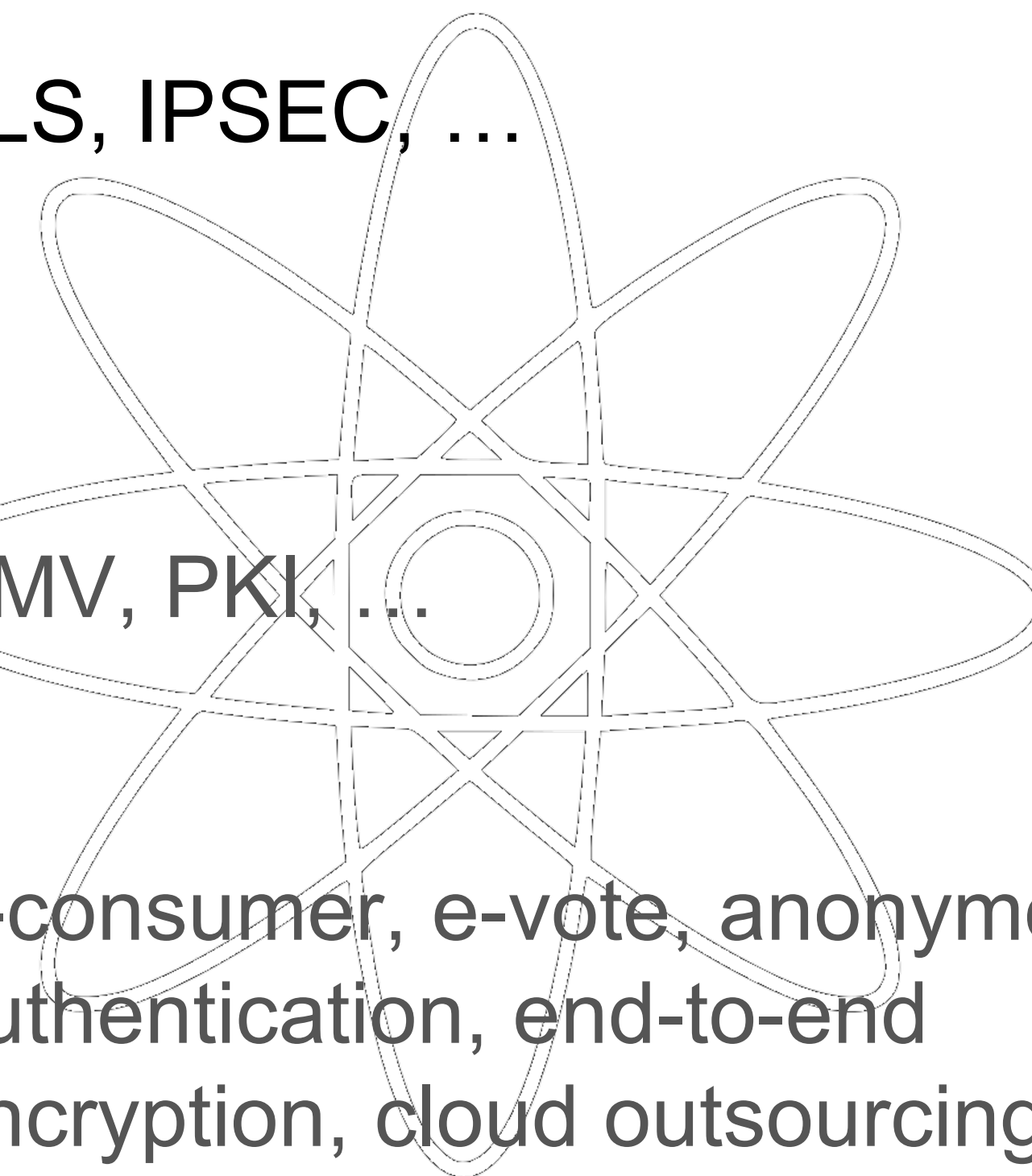
TLS, IPSEC, ...



EMV, PKI, ...

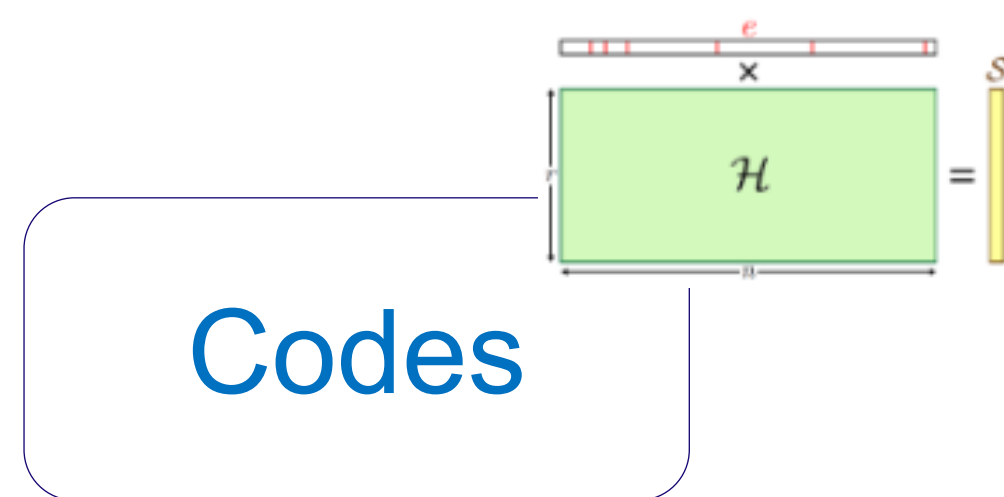


e-consumer, e-vote, anonymous authentication, end-to-end encryption, cloud outsourcing, ...



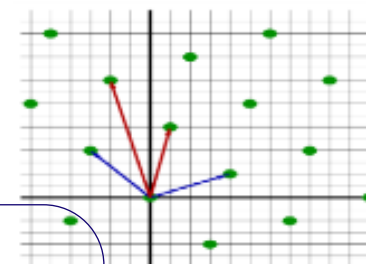
Post-Quantum Cryptography

- Post-Quantum Cryptography is related to **new mathematical problems** for which **quantum computers are not better** than classical ones
- Several practical solutions are known exist since mid 70s



Codes

Euclidean lattices



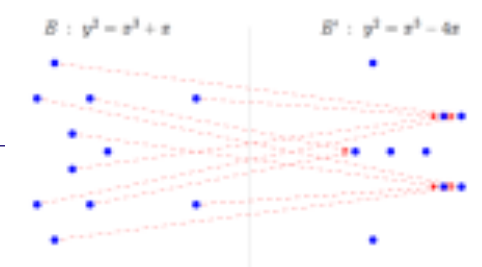
Multivariate polynomials

$$\begin{aligned}
 p^{(1)}(x_1, \dots, x_n) &= \sum_{i=1}^n \sum_{j=1}^m \beta_{ij}^{(1)} \cdot x_i x_j + \sum_{i=1}^n \beta_i^{(1)} \cdot x_i + \beta_0^{(1)} \\
 p^{(2)}(x_1, \dots, x_n) &= \sum_{i=1}^n \sum_{j=1}^m \beta_{ij}^{(2)} \cdot x_i x_j + \sum_{i=1}^n \beta_i^{(2)} \cdot x_i + \beta_0^{(2)} \\
 &\vdots \\
 p^{(m)}(x_1, \dots, x_n) &= \sum_{i=1}^n \sum_{j=1}^m \beta_{ij}^{(m)} \cdot x_i x_j + \sum_{i=1}^n \beta_i^{(m)} \cdot x_i + \beta_0^{(m)}
 \end{aligned}$$



Hash trees

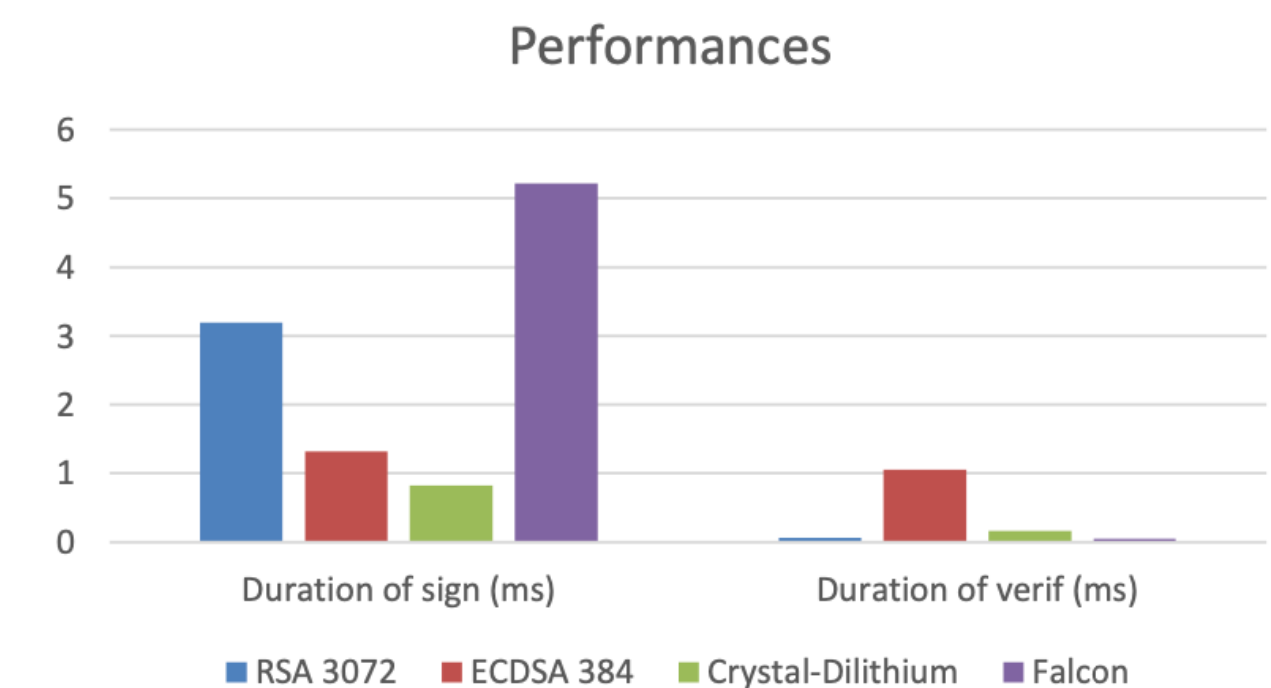
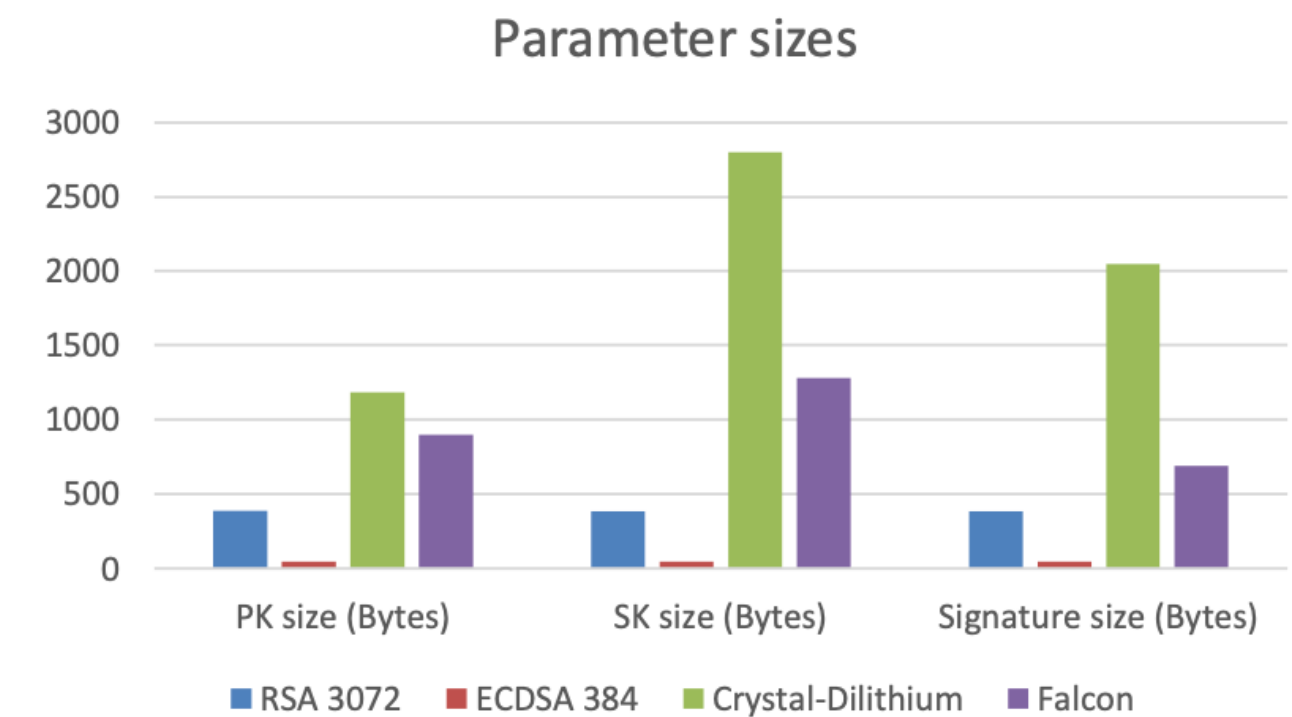
Isogenies



NIST standardisation process on PQC

Whole process

April 2015	Announcement	NIST announces future standardisation		
December 2017	Submission	69 complete and proper submissions	49 PKE/KEM	20 SIG
January 2019	End 1st round	26 submissions	17 PKE/KEM	9 SIG
July 2020	End 2 nd round	7 finalists, 8 alternates	4 (+5) PKE/KEM	3 (+3) SIG
July 2022	Winners	4 schemes selected for standardisation	1 PKE/KEM	3 SIG

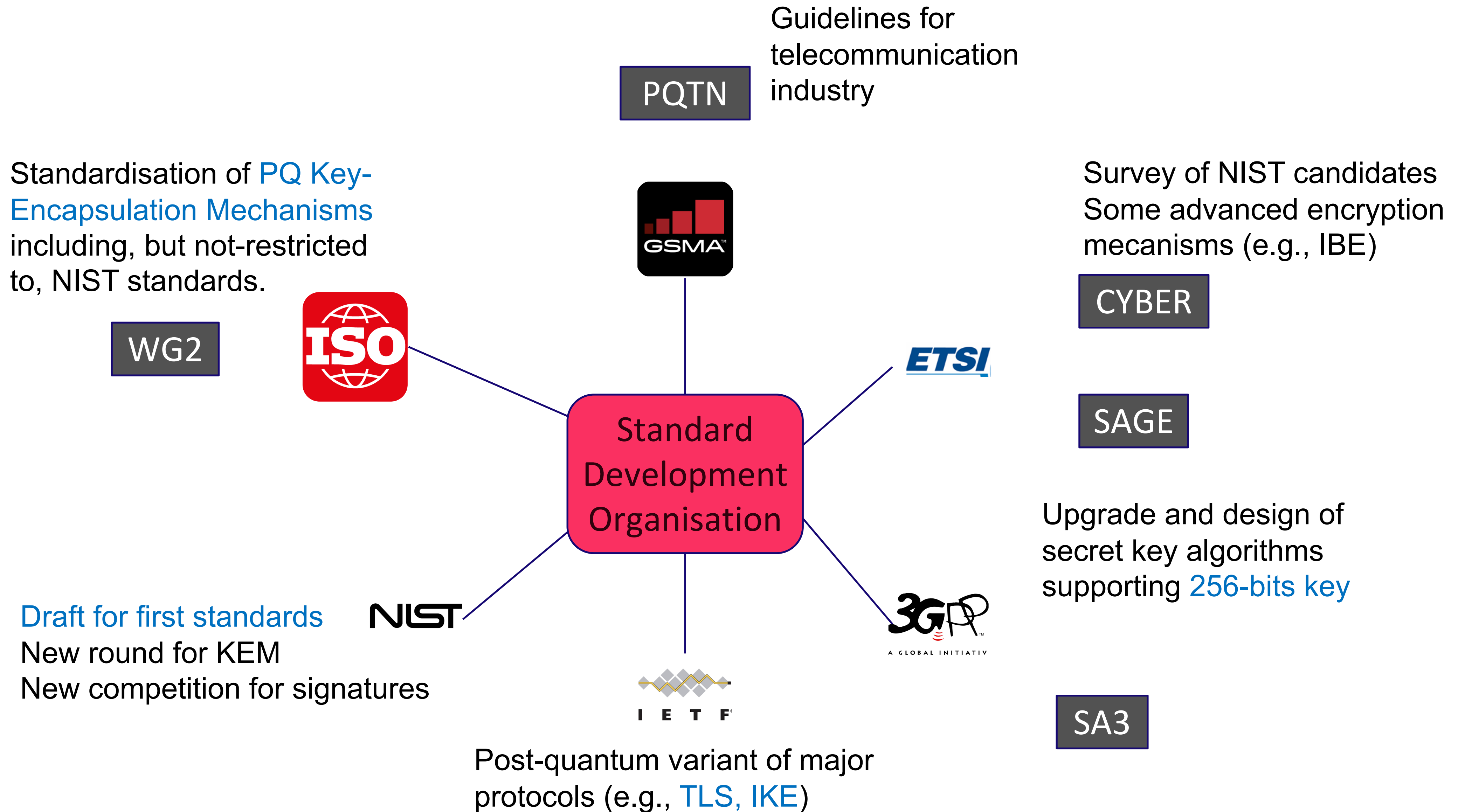


Selected candidates

PKE/KEM	Kyber (lattices)		
SIG	Dilithium (lattices)	Falcon (lattices)	SPHINCS+ (hash functions)

- + New round (PKE/KEM)
- + New competition (SIG)

Standardisation of PQC



What about advanced cryptographic mechanisms

Advanced encryption mechanisms		Advanced authentication mechanisms	
Privacy-preserving data treatment		Privacy-preserving authentication	Privacy-preserving payment
Fully homomorphic encryption	Other encryption with special features (e.g., IBE, ABE, FE)	Anonymous credentials/ attestations	E-cash
Inherently post-quantum secure (See next talk)	Hard to design post-quantum analogues at this stage	All those primitives are very good in the classical setting, leading to real products, or ready-to-market ones Not so easy in the post-quantum setting	

Our focus now

Anonymous credentials/attestations

- Advanced authentication mechanisms **enable full leakage control**
 - Blind Signature
 - Group Signature, DAA (Direct Anonymous Attestation), Enhanced Privacy ID (EPID)
 - Anonymous Credential
- Mechanisms **widely deployed** in billions of chips
 - Trusted Platform Module
 - Intel SGX enclaves
 - Willingness of the European Commission to deploy a privacy-preserving ID card...
- **Standardised** mechanisms
 - Blind signatures (ISO/IEC 18370)
 - Group signatures (ISO/IEC 20008)



Blind signature schemes



Message: m

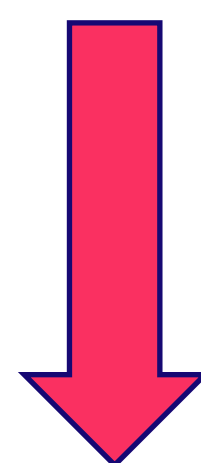
Mask the message



Signing key: sk

Sign the masked message from Alice

Unmask the signature $\Rightarrow \sigma$



(m, σ)

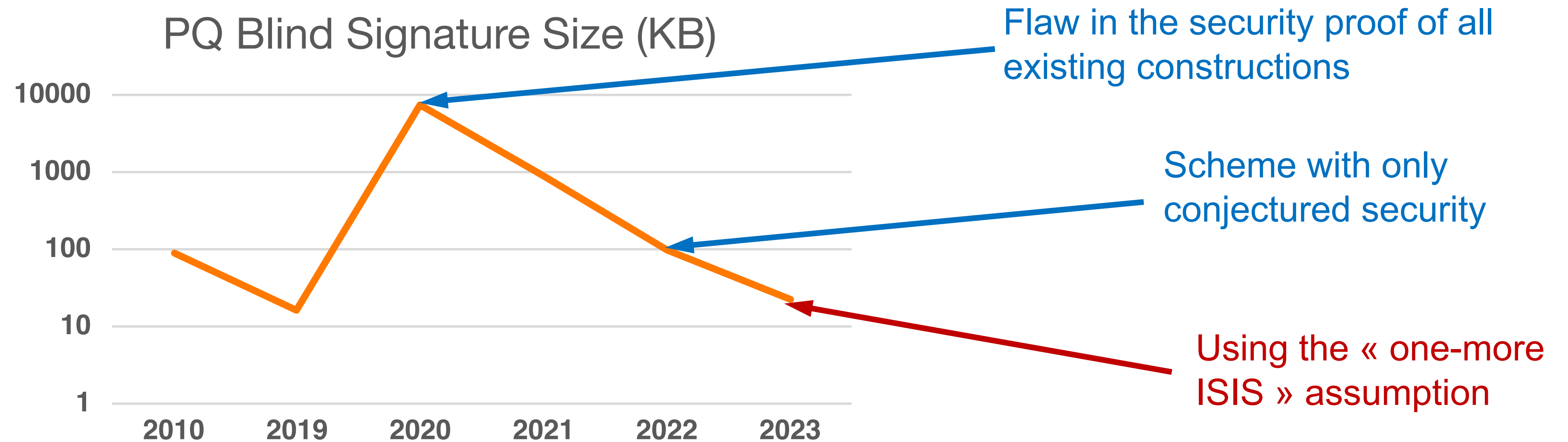


Verify validity of σ

- **Unforgeability**: infeasibility for Alice to create $l + 1$ message signature pairs after l interactions
- **Blindness**: infeasibility for signer to recognize the message signature pair

Post-quantum blind signatures

- **Standardised blind signatures** using classical cryptography
 - Based on RSA or Schnorr signatures
 - Size: 65 B
- Status on post-quantum constructions



Group signature schemes



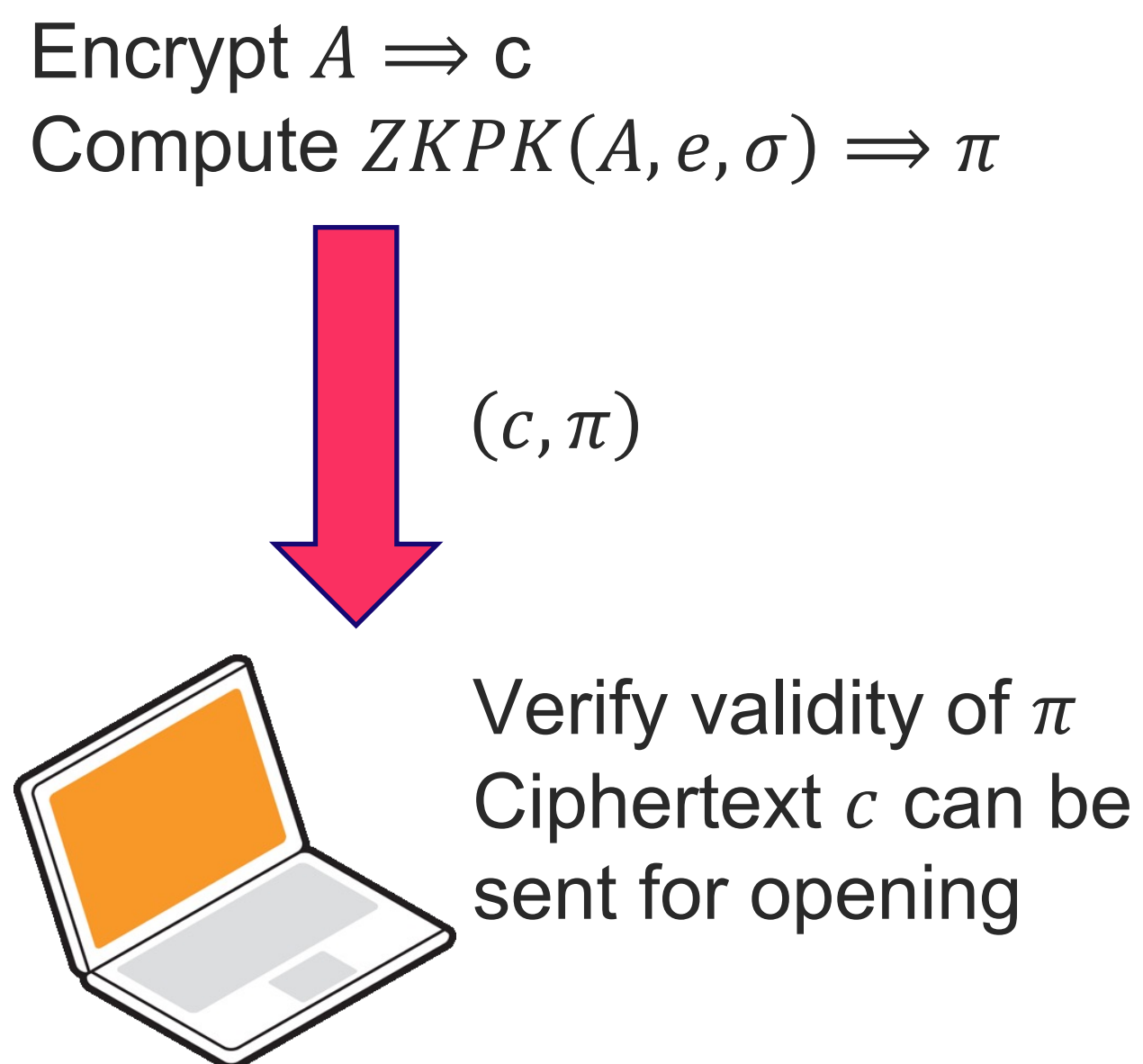
Secret: x
Public ID: A

Obtain a signature
on (x, A)



Signing
key: sk

Sign (x, A)



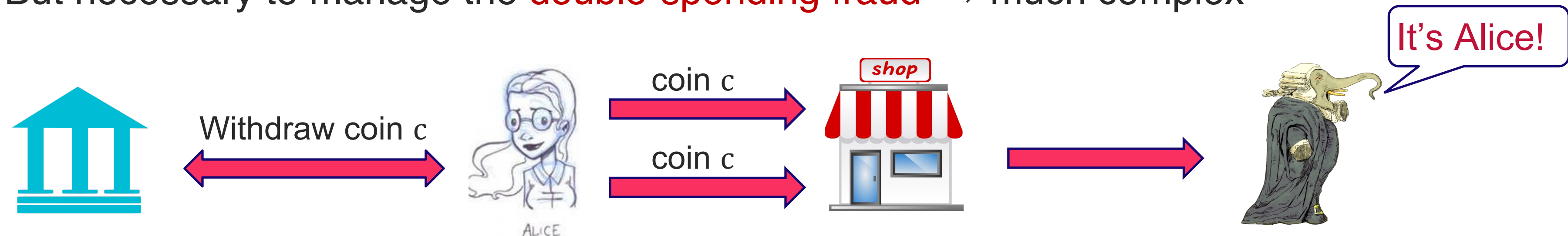
- **Unforgeability**: infeasibility to sign a message for non group members
- **Unlinkability**: infeasibility link two signatures from the same member
- **Traceability**: the opening of a valid signature should give the right member
- **Non-frameability**: infeasibility to falsely accuse a honest member

Post-quantum group signatures

- Basic tools to construct a group signature scheme
 - Signature scheme with advanced features + Encryption mechanisms
 - Zero-knowledge proof of knowledge (ZKPK)
- Standardised group signatures using classical cryptography
 - Based on pairings or on flexible RSA \Rightarrow Size: 160 B to 1 KB
- Status on post-quantum constructions
 - Hard to manage ZKPK efficiently compatible with signatures and encryption
 - Using a standard lattice assumption \Rightarrow Size: 600 KB
 - Using an interactive (stronger) version \Rightarrow Size: 30 KB

E-cash

- Main idea
 - Using group signatures: **a wallet is a group member**
 - But necessary to manage the **double-spending fraud** \Rightarrow much complex



- E-cash constructions using classical cryptography
 - Using pairings, implemented in a smart phone: **payment < 100 ms**
- Status on post-quantum constructions
 - Security proof of most of the constructions has been invalidated
 - One generic scheme but no instantiation on post-quantum cryptography
 - One concrete scheme using lattices \Rightarrow **Factor of 1 million compare to pairing-based solutions!**

Conclusion

Quantum computers are not for now
But the threat is huge, and we need to be prepare

- **Basic cryptography**

- NIST standards are in progress, other standards are working
- Integration is the next steps \Rightarrow should be done by the industry
- But research should continue
 - Cryptanalysis
 - Improve efficiency
 - Hardware implementation and security (SCA)

- **Advanced cryptography**

- We are very far from what can be done using classical cryptography
- One of the main research topic on privacy-preserving cryptography
 - Pairings look impossible
 - Use stronger assumptions to improve efficiency
 - Necessity to think differently

Thank you

