

Définition Formelle de la Relation de Dépendance Causale entre Événements Journalisés

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CentraleSupélec & Thales

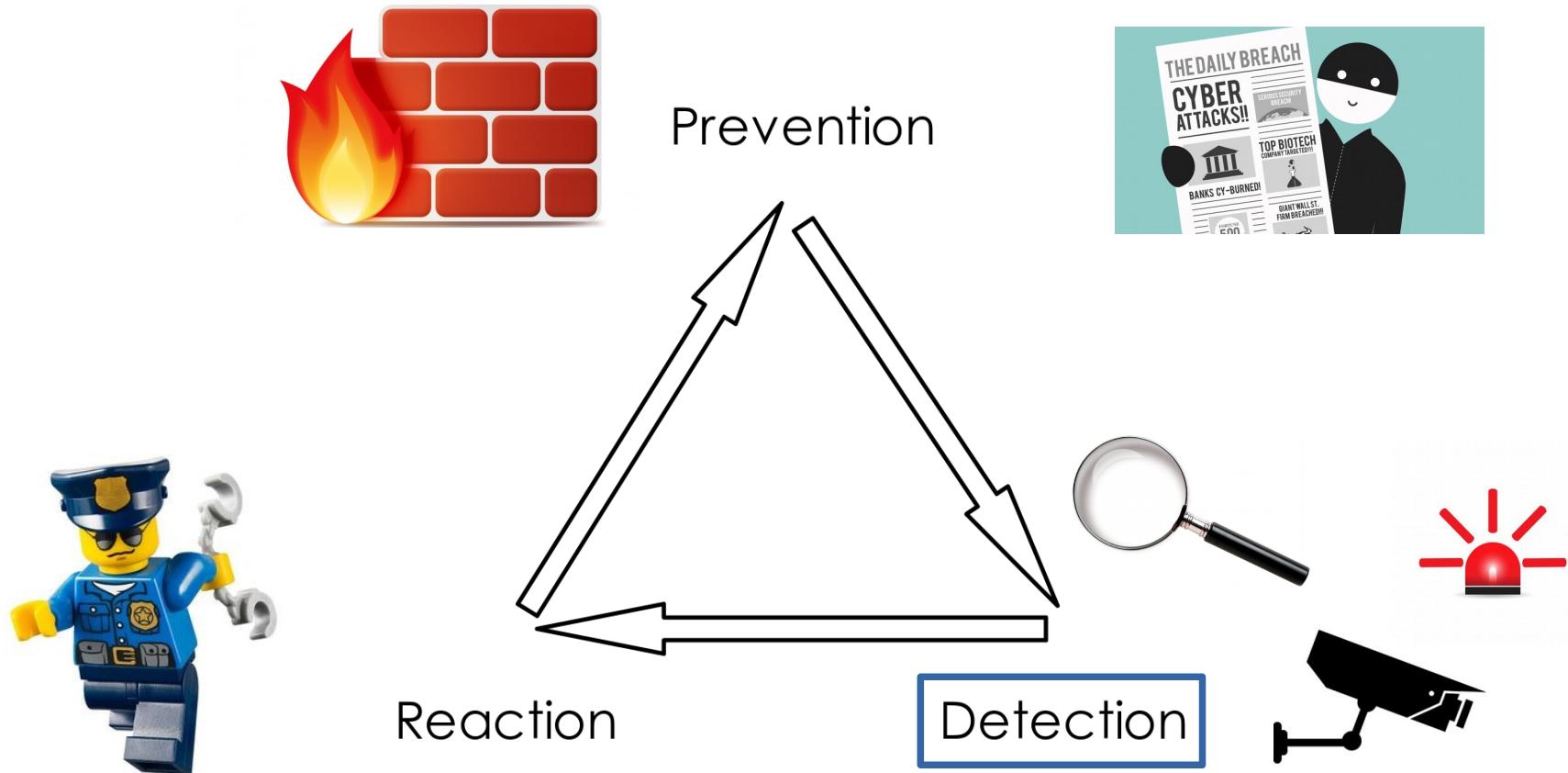
Encadrants :
Eric TOTEL - CentraleSupélec
Olivier BETTAN - Thales

Agenda

Objectif : Susciter votre Curiosité pour la session Poster

- Contexte
- Objectif
- Contribution

Contexte – Un autre triangle de la Sécurité

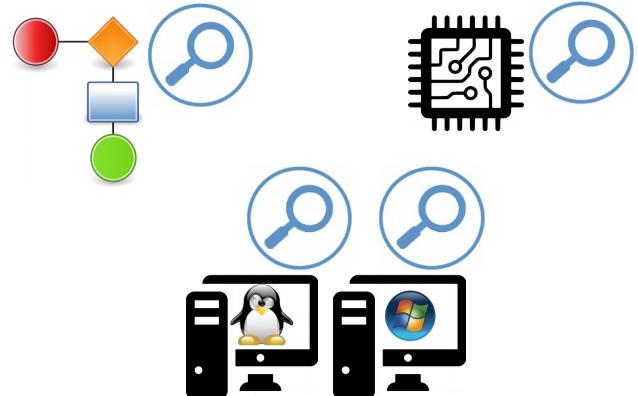


Contexte – Le besoin de Supervision en Sécurité

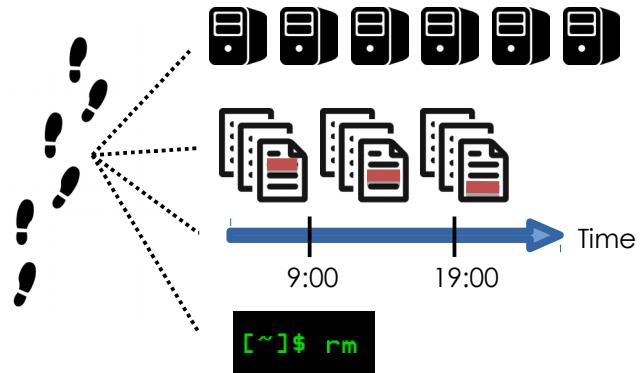
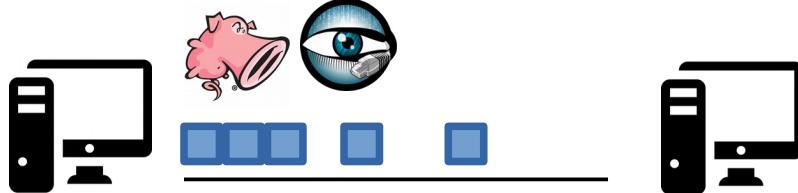


Différentes couches d'abstraction :

- Application (App Logs, CPU, ...)
- Système d'exploitation (Syscalls, ...)
- Réseau (DPI, NIDS, Netflow, ...)

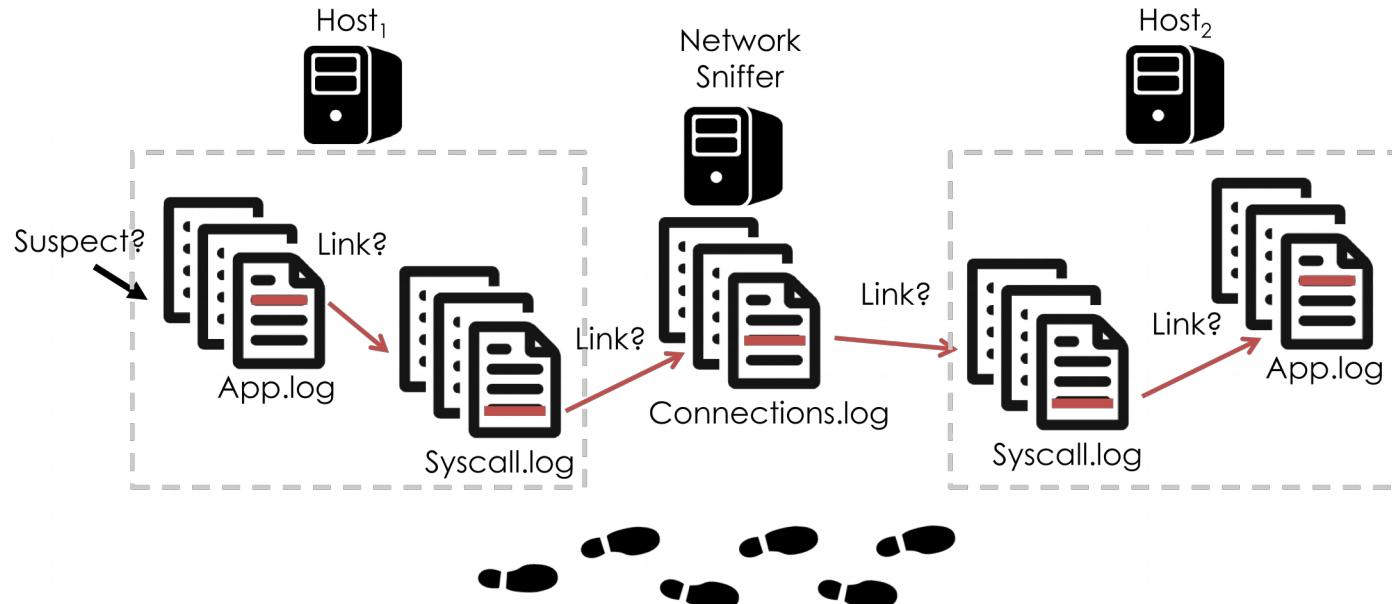


Observation partielle des actions effectuées



Objectif – Ce qu'un Analyste fait

Recherche de liens de corrélation entre les événements journalisés

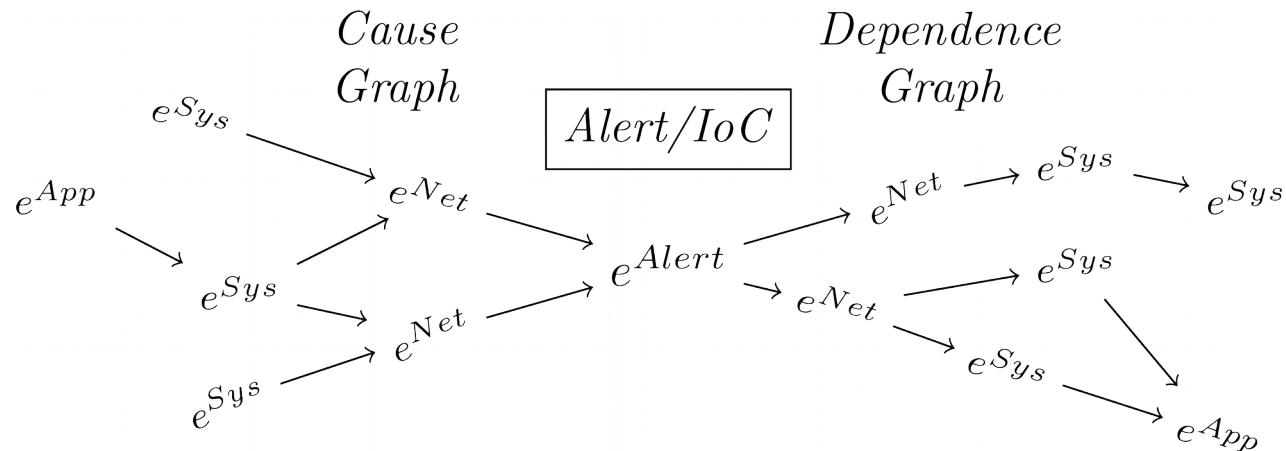


Objectif – Ce qu'un Analyste veut

Retrouver l'ensemble des **actions** effectuées par un attaquant

⇒ Découverte de scénarios d'attaque

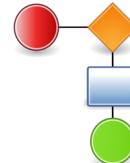
⇒ Recherche des liens de **dépendances causales** entre les événements correspondant aux actions de l'attaquant



Objectif – Définition de la Dépendance Causale

Constat :

Pas de définition formelle de la dépendance causale entre événements journalisés hétérogènes



*definition of
causal dependency?*

$$e_1 \xrightarrow{?} e_2$$



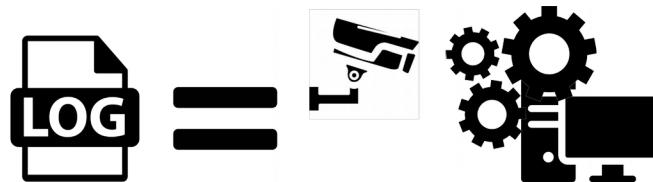
Modélisation d'Attaque (Arbres d'Attaque, Graphes d'Attaque, ...)

Contribution – Raisonnement

Objectif : Définition de la relation de Dépendance Causale entre Événements

Définition d'Événement : « une action identifiable ayant lieu sur un dispositif et étant enregistrée comme une entrée de journal ».

[European Commission, 2010]



Contribution – À propos d'Actions et d'États

Relation de D'Ausbourg [d'Ausbourg, 1994] entre les **états** des **objets** d'un système

⇒ flux d'information

$$(o, t) \rightarrow (o', t')$$

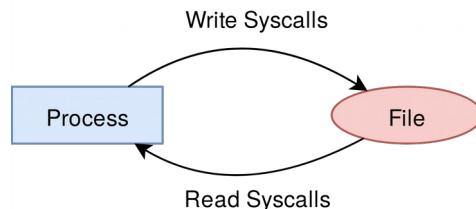
$$a := 1$$

$$a := a + 1$$

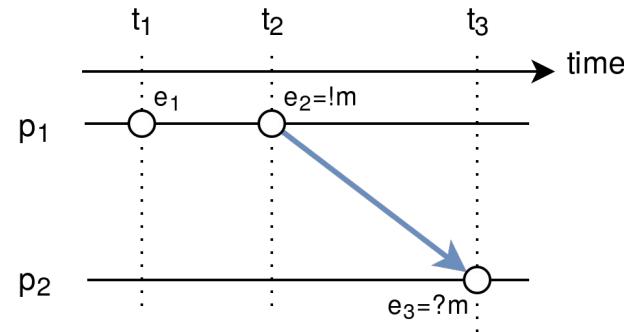
$$a := 0$$

$$(a, 1) \rightarrow (a, 2)$$

$$(a, 2) \vdash \rightarrow (a, 3)$$



Relation de Lamport [Lamport, 1978] entre les **actions** effectuées par les processus d'un système distribué.



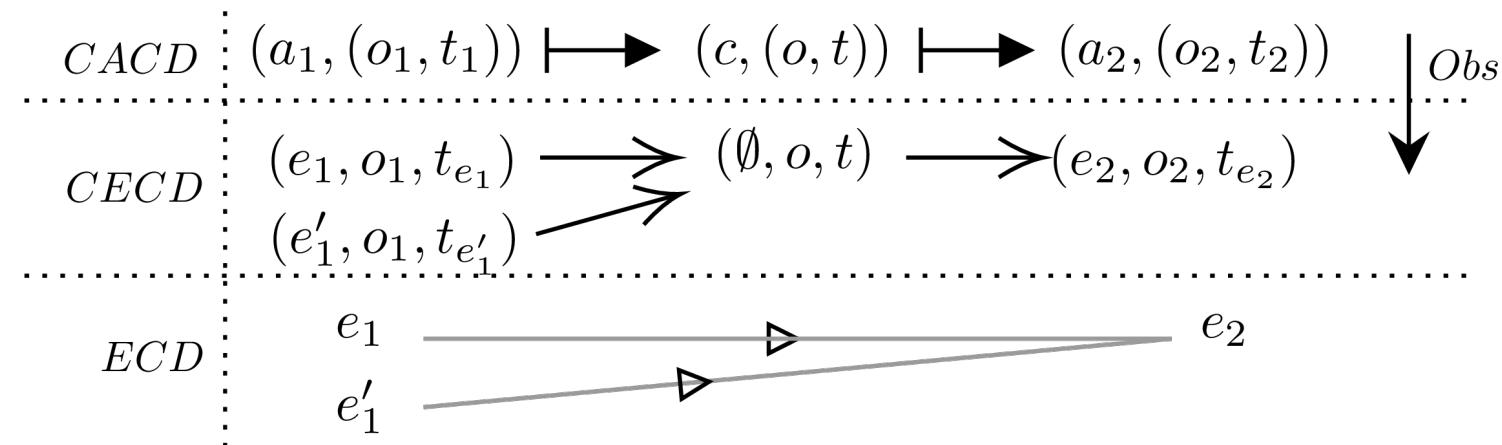
[d'Ausbourg, 1994] Implementing secure dependencies over a network by designing a distributed security subsystem. ESORICS.

[Lamport, 1978] Time, clocks, and the ordering of events in a distributed system. Communications of the ACM.

Contribution – Nouvelles Dépendances Causales

Définitions de 3 nouvelles relations de dépendances causales :

1. Actions Contextuelles (CACD)
2. Événements Contextuels (CECD)
3. Événements Bruts (ECD)



Conclusion

- Nouveau cadre théorique
 ⇒ Dépendance Causale entre Événements Journalisés ;
- Découverte de liens entre Événements ;
- Implémentation dans un environnement Linux

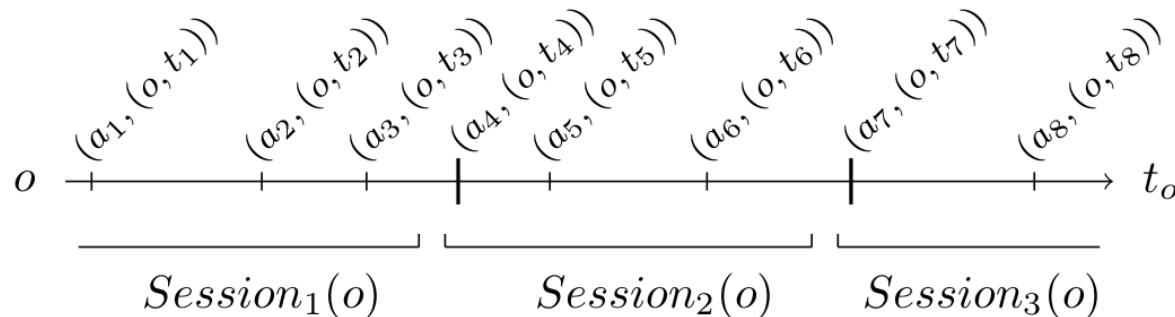
Merci de votre attention

Contribution – Introduction des Sessions

Volonté d'avoir un modèle plus précis

⇒ 2 États consécutifs peuvent être indépendants

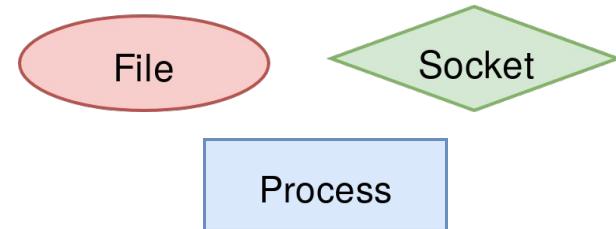
$$\begin{aligned} Session_n(o) = \{ & (a_i, (o, t_i)) / \\ & (o, t_i) \rightarrow (o, t_{i+1}) \\ & \wedge (o, t_{end_{n-1}}) \not\rightarrow (o, t_{start_n}) \\ & \wedge (o, t_{end_n}) \not\rightarrow (o, t_{start_{n+1}}) \} \end{aligned}$$



Contribution – [Def] Action Contextuelle

Deux types d'Objets :

- Passifs (ex : containers d'information) => Actions
- Actifs (ex : processus ou le réseau) => Actions



Une Action est effectuée dans un certain État

Def : Action Contextuelle ($a, (o, t)$)

- $(Action, \text{État}) == (Action, (\text{Objet}, \text{Horodatage}))$
- $a \in \text{ObjectActions}(o)$ avec $\text{ObjectActions}(o) = \{ a_i \} \cup \{\emptyset\}$

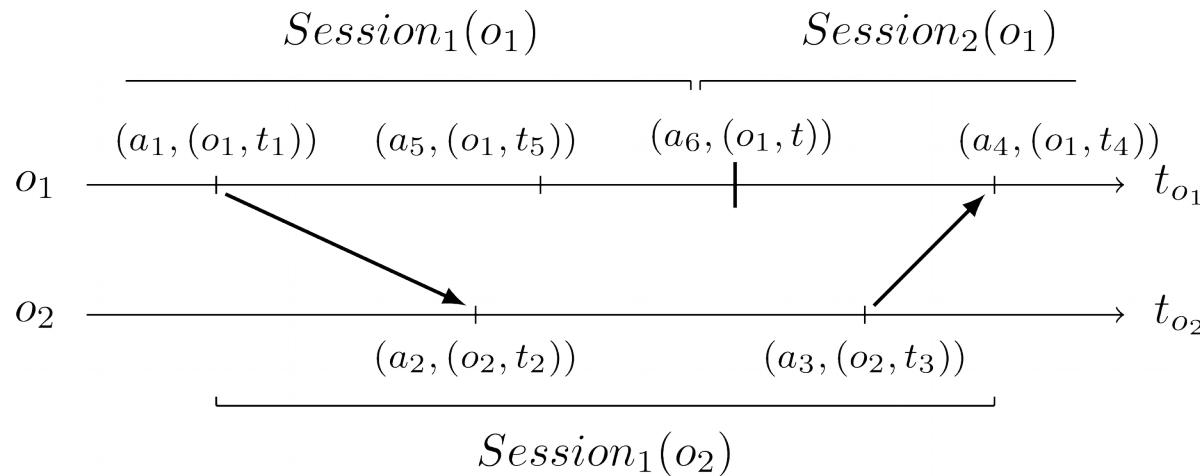
Contribution – Action Contextuelle

Dépendances Causales :

- Actions effectuées par les Objets (ex : send & recv message)
- États des Objets (c-à-d flux d'information)

(Action, (Objet, Horodatage))

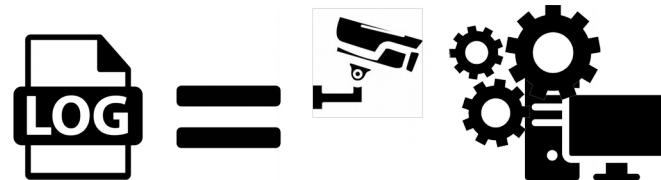
État



Contribution – [Def] Événement Contextuel

Action Contextuelle ($a, (o, t)$) :

- Observée ... ou pas
- Par une ou plusieurs sondes

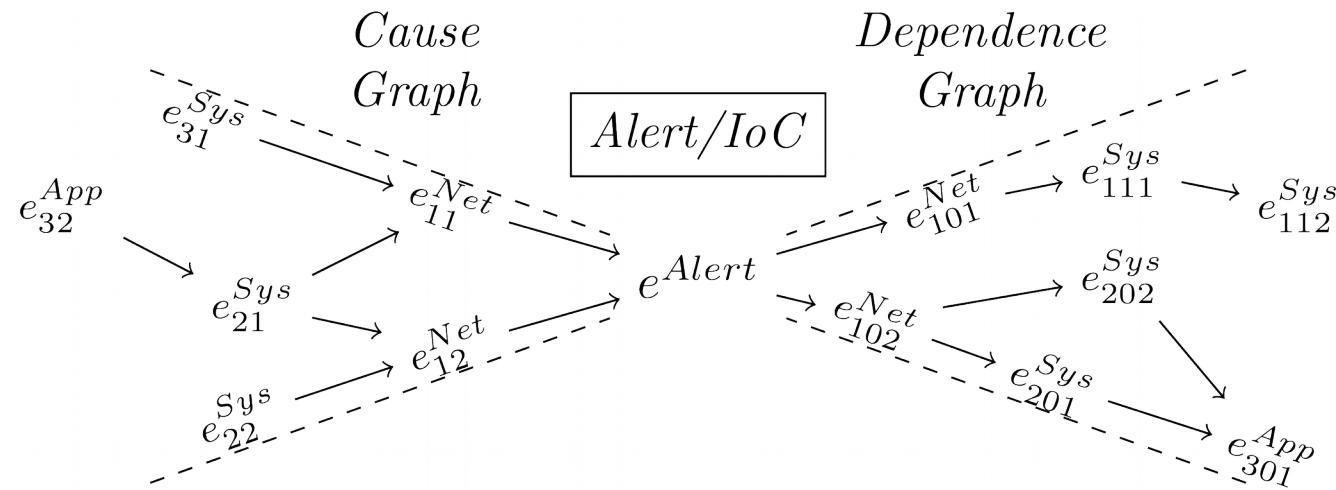


Événement Contextuel (e, o, t_e) \Rightarrow Observation d'une Action Contextuelle

$$\text{Obs}((a, (o, t_a))) = \{(e_i, o, t_{ei})\} \cup \{(\emptyset, o, t_a)\}$$

Contribution – Cause & Dependence Graphs

Cause & Dependence Graphs can be computed for each layer depending on the use-case

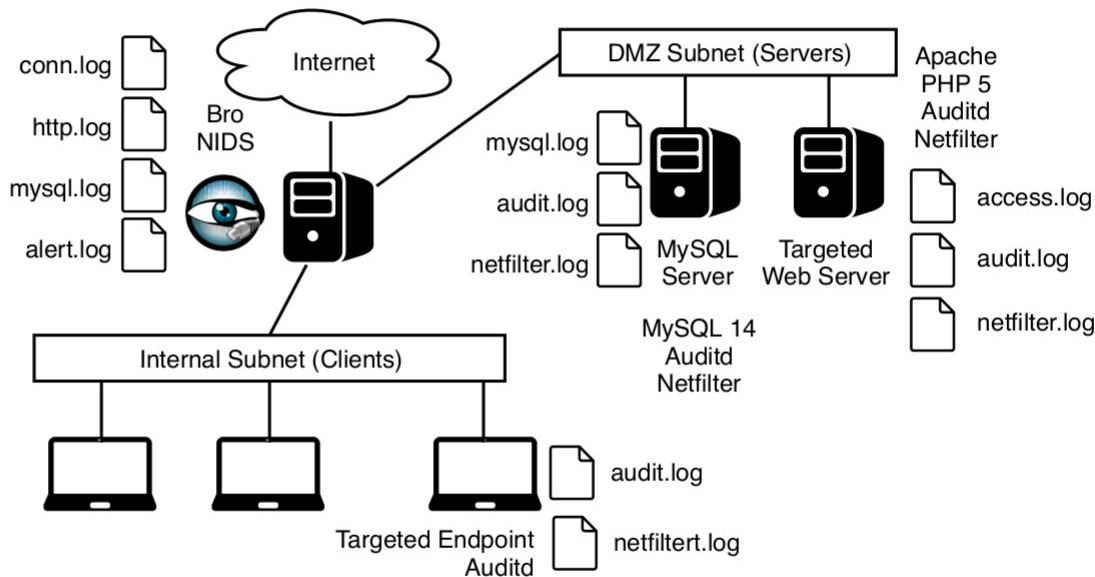


Evaluation – Challenges

- No datasets with heterogeneous events publicly available
- Need to create our own test environment
- Need to elaborate our own Attack Scenarios
- How to evaluate our implementation of the model ?

Evaluation – Test Environment

Leveraging Logged Events Semantics to Compute an Approximation of the Contextual Event Causal Dependency Layer.

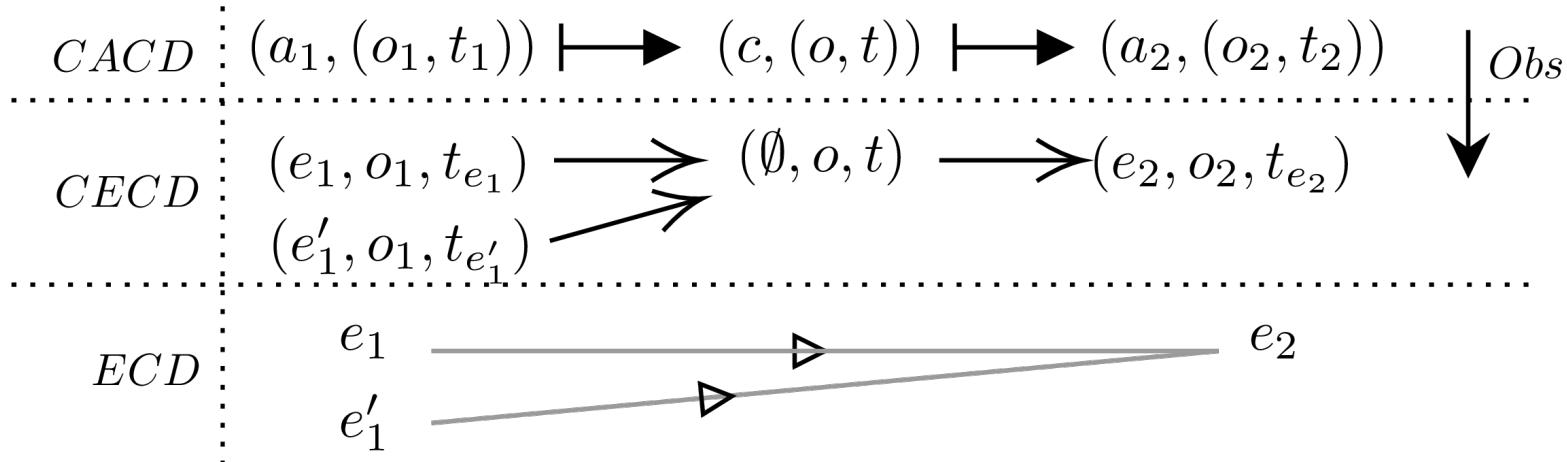


Evaluation – Attack Scenarios & Data Generation

1. SQL Injection against a vulnerable PHP Script;
2. Trojan Software against an End Point machine (Ubuntu);
3. Command Injection Attack leveraging the ShellShock Bash Vulnerability (CVE-2014-6271) against Apache Web Server.

Evaluation – Two Point of Views

- Bottom-Up Approach => Current Practices in SIEM
- Top-Down Approach => Instrumentation



Future Works

- Applying the bottom-up to Windows use-cases
Leveraging Syscalls API using ProcMon or LogMan
- Top-down approach by building all layers from Contextual Actions to Contextual Events.
Towards a record and replay system to compute objects' states.
==> Leveraging Dynamic Information Flow Tracking (DIFT)
Message passing system within the Network layer of the Kernel.

Conclusion

- Introduction d'un nouveau cadre théorique permettant de raisonner sur la notion de Dépendance Causale entre Événements Journalisés ;
- Bottom-up approach with a lightweight implementation.
Building an approximation of the model from the logged events.
- Current methods and implementations allow the observation and recording of different subsets of actions.
/\ existing work only enables an approximation of the correct model.

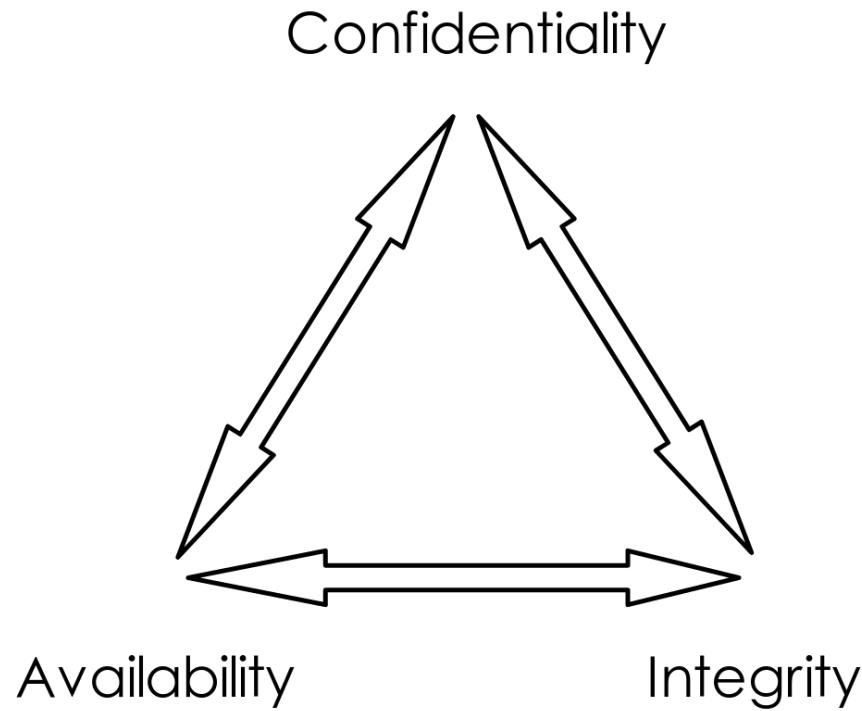
Merci de votre attention

Qui je suis

Charles XOSANAVONGSA

- Doctorant CIFRE en 3^{ème} année
- CentraleSupélec – Équipe CIDRE (Rennes)
- Thales Six GTS France – Équipe Études Amonts (Palaiseau)
- Sécurité Informatique
 - Détection d'intrusion
 - Analyse de Logs
 - Corrélation d'alertes & d'événements

Context – Security Triangle



Context – You will be breached...

Attacker's Goal:

Gaining Foothold

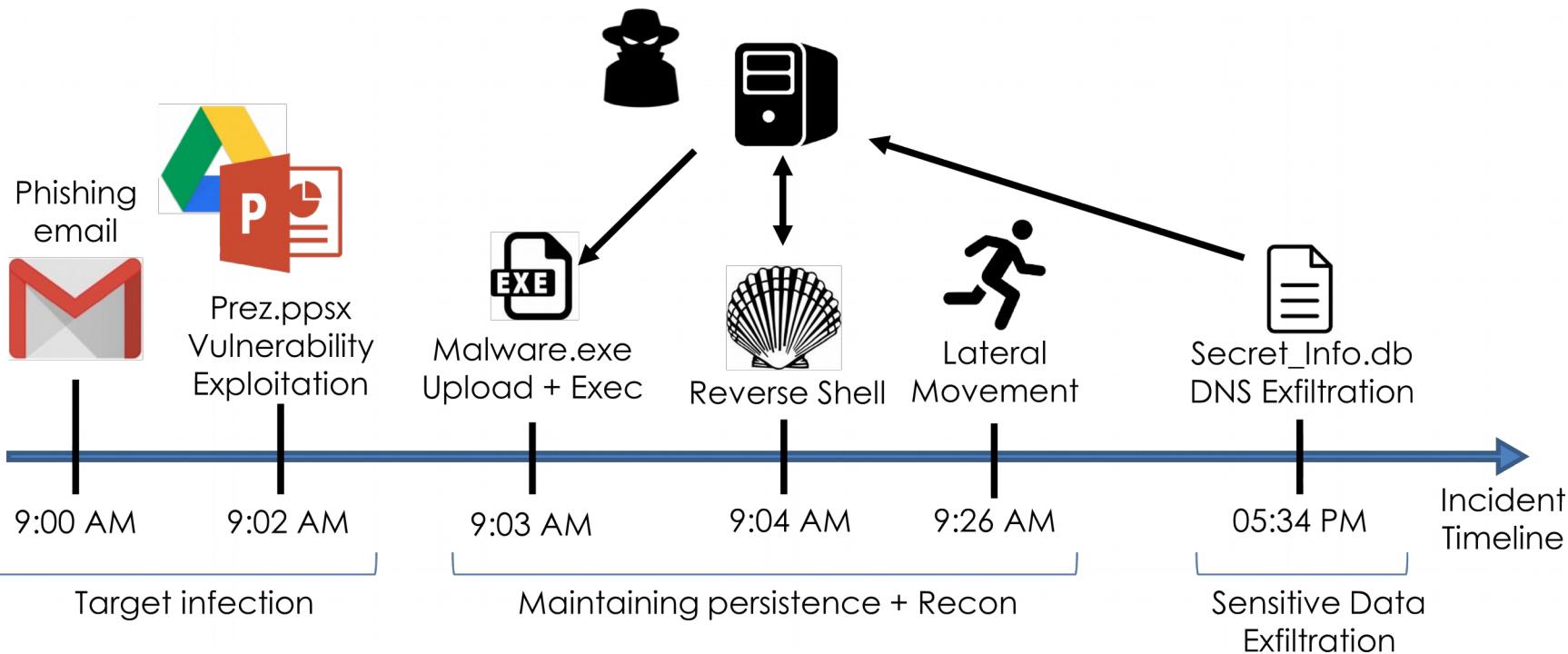
Observation:

They eventually succeed

~~Prevention Mechanisms~~



Context – Attack Scenario Example



Related Work – Journey Overview

- Alert & Event Correlation
- Information Flow Tracking
- D'Ausbourg's Causal Dependency among Objects' States
- Lamport's Happened-Before Relation among Processes' Actions

Related Work – Explicit Alert & Event Correlation

- Attack Specification based
- Expression of Cause and Effects Relations between Events
- What we want:
Discovering Attack Scenarios, without specifying them,
through Heterogeneous Logs Analysis

Related Work – D'Ausbourg's Model

Relation of **causal dependency**, leveraging information flows between states of the system : $(o,t) \rightarrow (o',t')$

A state (o, t) is the value of an **object** o at a given time t .

No constraints on the definition of objects

- Program Variables
- Files, Sockets, Pipes, Memory, Processes...

$a := 1$

$a := a + 1$

$a := 0$

$(a,1) \rightarrow (a,2)$

Drawbacks :

- Does not take objects' actions into account

$(a,2) \leftarrow (a,3)$

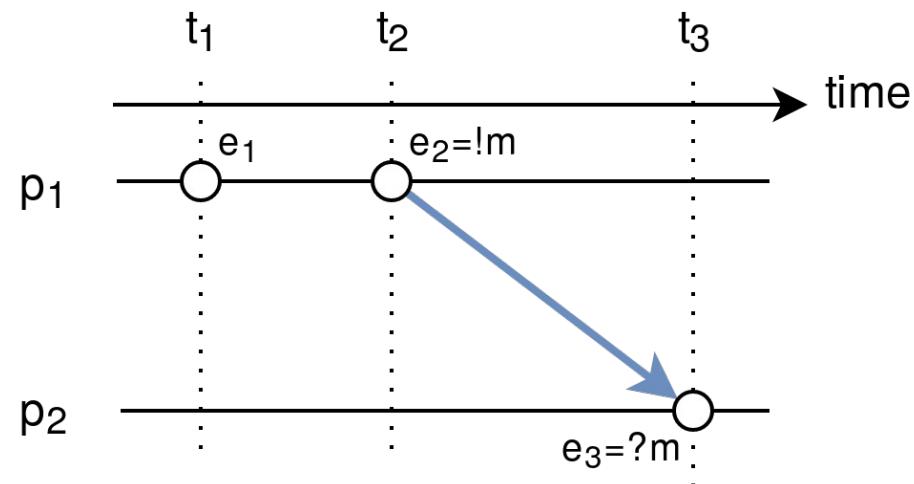
Related Work – Lamport's Model

Temporal causality between **actions** performed by processes of a **distributed system**.

=> No **global clock** in the distributed system.

Partial order relation. $a < b$ if :

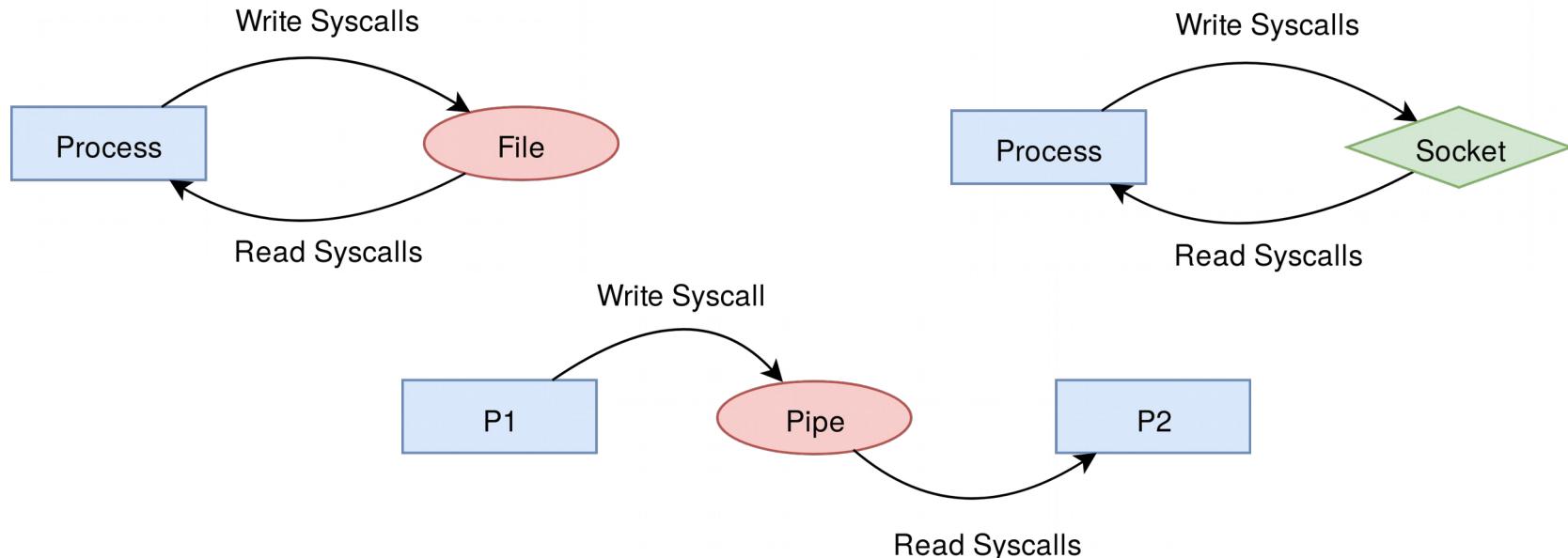
1. a is performed before b on the same process
2. $a = !m \&& b = ?m$
3. $a < c \&& c < b$



Related Work – Information Flow Tracking (IFT)

Monitoring Information Flows between the System's Objects

Use-case: Security Policies for Confidentiality and Integrity



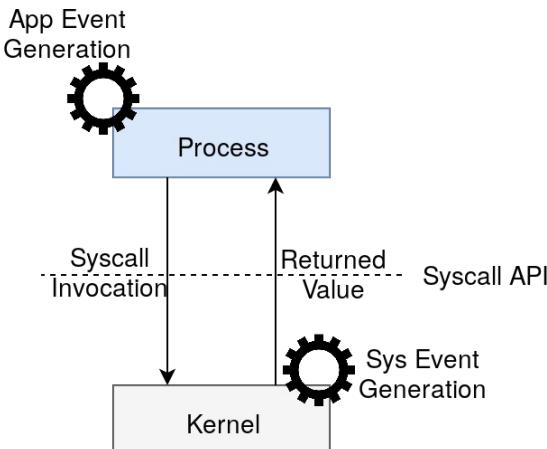
Related Work – Information Flow Tracking (IFT)

Allows to answer the following questions:

- Where does a given object come from?
- What are the objects that it influences?

Can be implemented in different Abstraction Layers:

- Dynamic Data Flow Analysis (CPU)
[Kemerlis et al. 2012]
- Syscall (OS) [Georget et al. 2017]
- Application + OS [Muniswamy-Reddy et al. 2009]



[Kemerlis et al. 2012] libdft: Practical Dynamic Data Flow Tracking for Commodity Systems.

[Georget et al. 2017] Information Flow Tracking for Linux Handling Concurrent System Calls and Shared Memory.

[Muniswamy-Reddy et al. 2009] Layering in Provenance Systems.

Related Work – Dependency Explosion

Reasoning on Syscalls is too coarse-grained:

Over Approximation of Causal Dependencies

Given a Process, a given Syscall is supposed to be dependent on all the previous Syscalls invoked.

Prior work that proposes a solution:

- Binary Analysis in order to Identify the different Units inside processes and the information flows among them. [Lee et al. 2013]
=> Leverages LibC & Binary Instrumentation
- Rareness Score Computation in order to identify paths that might represent attack scenarios. [Liu et al. 2018]

[Lee et al. 2013] High Accuracy Attack Provenance via Binary-based Execution Partition.

[Liu et al. 2018] Towards a Timely Causality Analysis for Enterprise Security.

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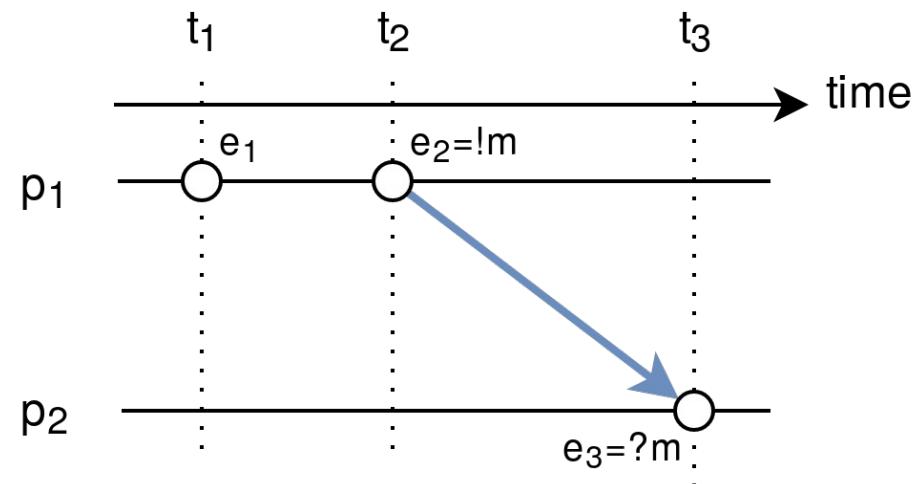
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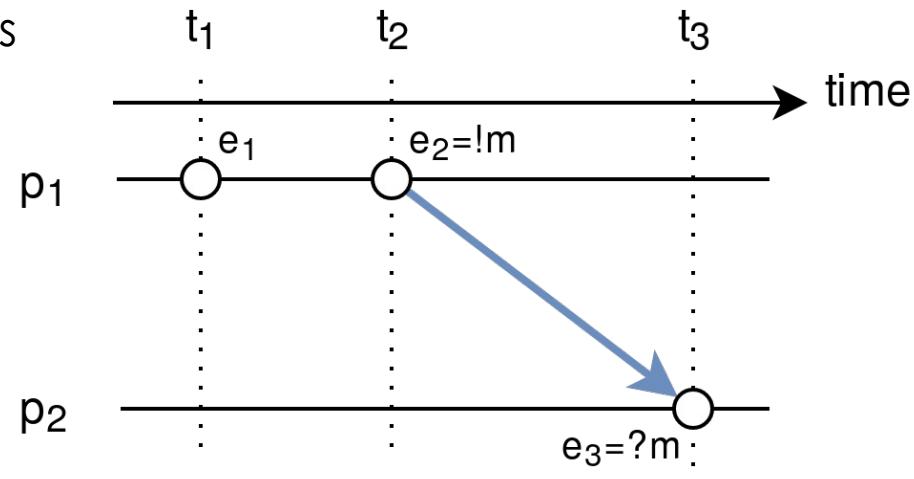
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2. $a = !m \&& b = ?m$
3. $a < c \&& c < b$



Related Work – Lamport's Model

Drawbacks :

- Over approximation of Causal Dependencies
- Only deals with application level actions
=> No heterogeneous events

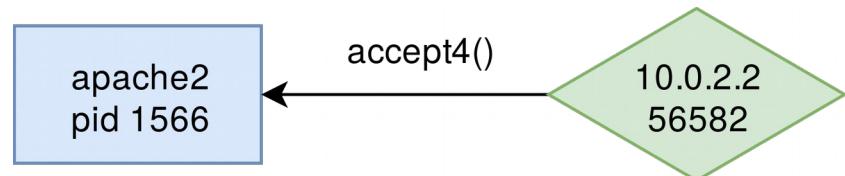


Evaluation – Bottom-Up Approach – Data Generation

Leveraging Logged Events Semantics to Compute an Approximation of the Contextual Event Causal Dependency Layer.

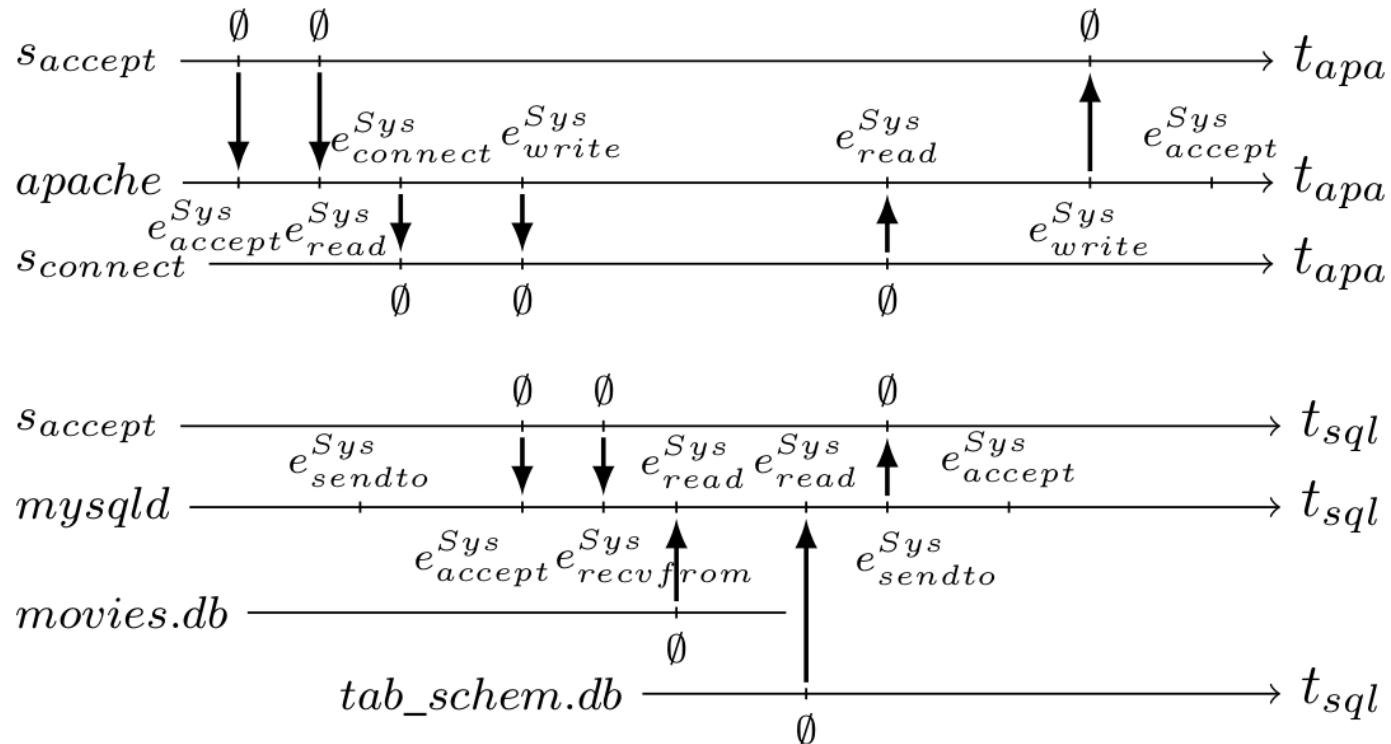
Syscall Netfilter PCAP Application

```
type=SYSCALL msg=audit(1541366508.539:47875): arch=c000003e syscall=288 success=yes  
exit=10 a0=3 a1=7ffce59a1100 a2=7ffce59a10e0 a3=80000 items=0 ppid=1106 pid=1566  
auid=4294967295 uid=33 gid=33 euid=33 suid=33 egid=33 sgid=33 fsgid=33  
tty=(none) ses=4294967295 comm="apache2" exe="/usr/sbin/apache2" key=(null)  
type=SOCKADDR msg=audit(1541366508.539:47875): saddr=0200DD060A0002020000000000000000  
(saddr= (AF INET) 10.0.2.2 : 56582)  
type=PROCTITLE msg=audit(1541366508.539:47875):  
proctitle=2F7573722F7362696E2F61706163686532002D6B007374617274  
(proctitle=/usr/sbin/apache2 -k start)
```



Evaluation – Illustration of Contextual Event

Syscall : 1 Active Object + 1 Passive Object + 1 Information Flow Action



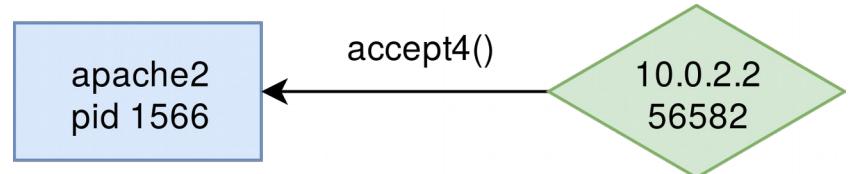
Evaluation – Bottom-Up Approach – Data Generation

Leveraging Logged Events Semantics to Compute an Approximation of the Contextual Event Causal Dependency Layer.

Syscall **Netfilter** PCAP Application

A connection is identified by a quadruplet
IP/Port Src & Dst

Bridging the gap between Syscalls and PCAP



Evaluation – Bottom-Up Approach – Data Generation

Leveraging Logged Events Semantics to Compute an Approximation of the Contextual Event Causal Dependency Layer.

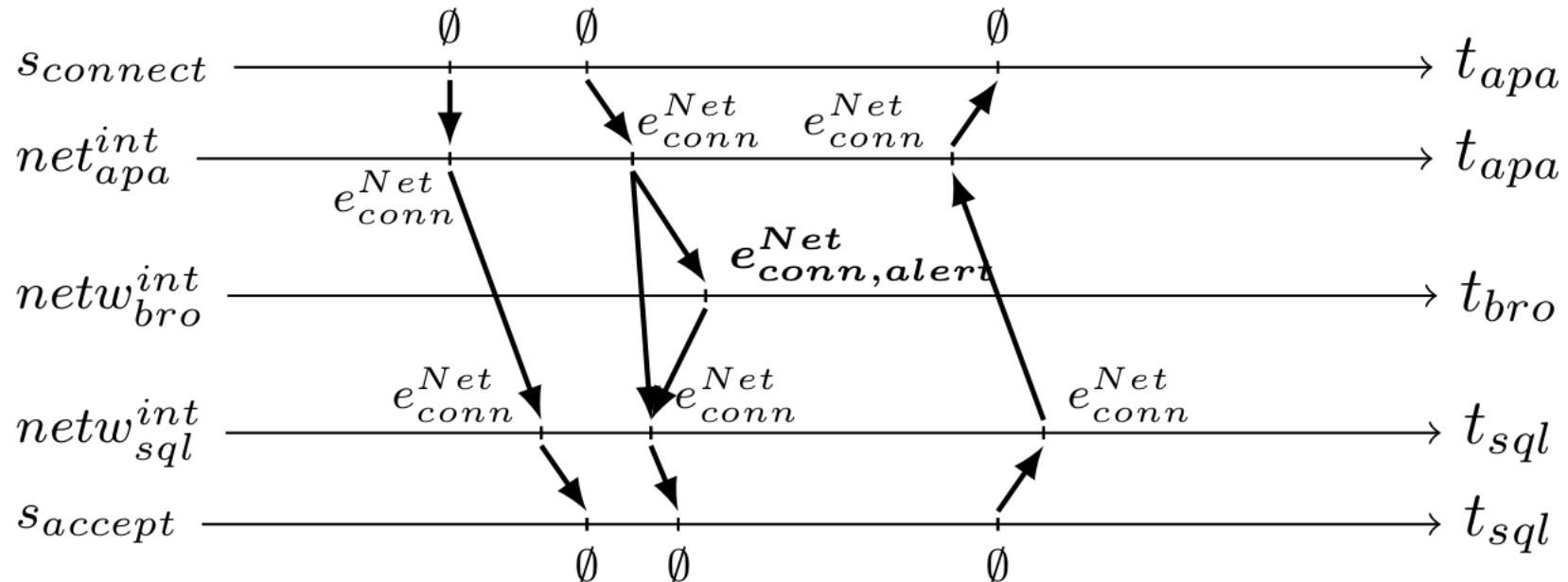
Syscall Netfilter PCAP Application

2 Passive Objects (Network Sockets) + Message Passing

2018-11-04T21:50:55.001600Z CgGkAp4P6ThQzD2Wg **192.168.1.2 48218 192.168.1.3 3306**
tcp MySQL::Sqlⁱ SELECT * FROM movies WHERE title LIKE ‘%’ UNION ALL SELECT
table schema,table name, null, null, null, null, null from information schema.tables;-
%’ SQLⁱ Attempt : Suspect syntax detected. [‘Notice::ACTION LOG’]

Evaluation – Illustration of Contextual Event

PCAP : 2 Passive Objects (Network Sockets) + Message Passing



Evaluation – Bottom-Up Approach – Data Generation

Leveraging Logged Events Semantics to Compute an Approximation of the Contextual Event Causal Dependency Layer.

Syscall Netfilter PCAP Application

1 Active Object (The process with its PID)

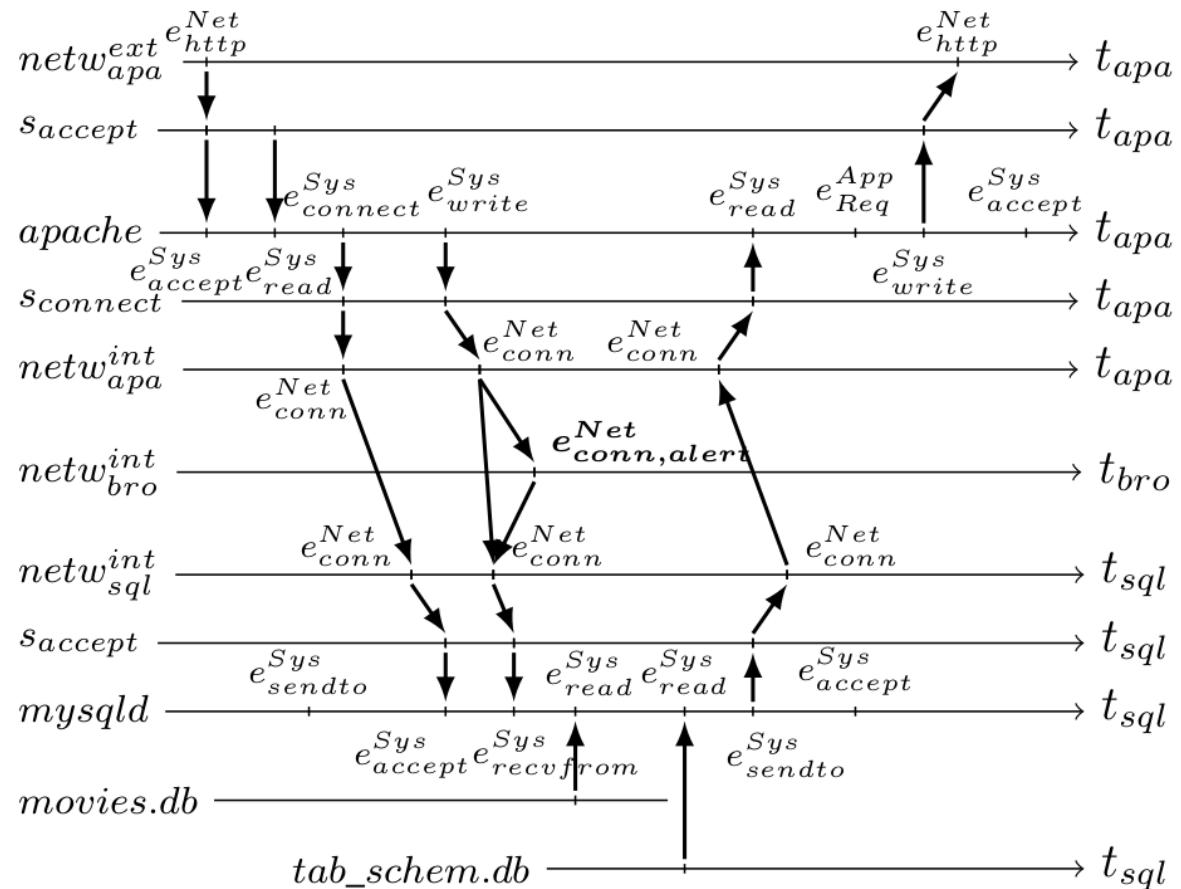
[04/Nov/2018:21:50:54 +0000] 1566 10.0.2.15 80 10.0.2.2 56582

“POST /bWAPP/sqli_6.php HTTP/1.1” 200 6799 “http://10.0.2.15:80/bWAPP/sqli_6.php”

“Mozilla/5.0 (X11; Ubuntu; Linux x86_64; rv:61.0) Gecko/20100101 Firefox/61.0”

Evaluation – Illustration of Contextual Event

Putting it all together



Evaluation – ShellShock attack against Apache

Leveraging Logged Events Semantics to Compute an Approximation of the Contextual Event Causal Dependency Layer.

